

# PostgreSQL Performance

## The basics



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# The dumb simple

RAID 1 or 10 (RAID 5 is for chumps)

Lots of memory. It's cheap, get 4G (minimum)

At least two cores (PostgreSQL is process based)

If SATA consider twice as many disks

Always get a BBU

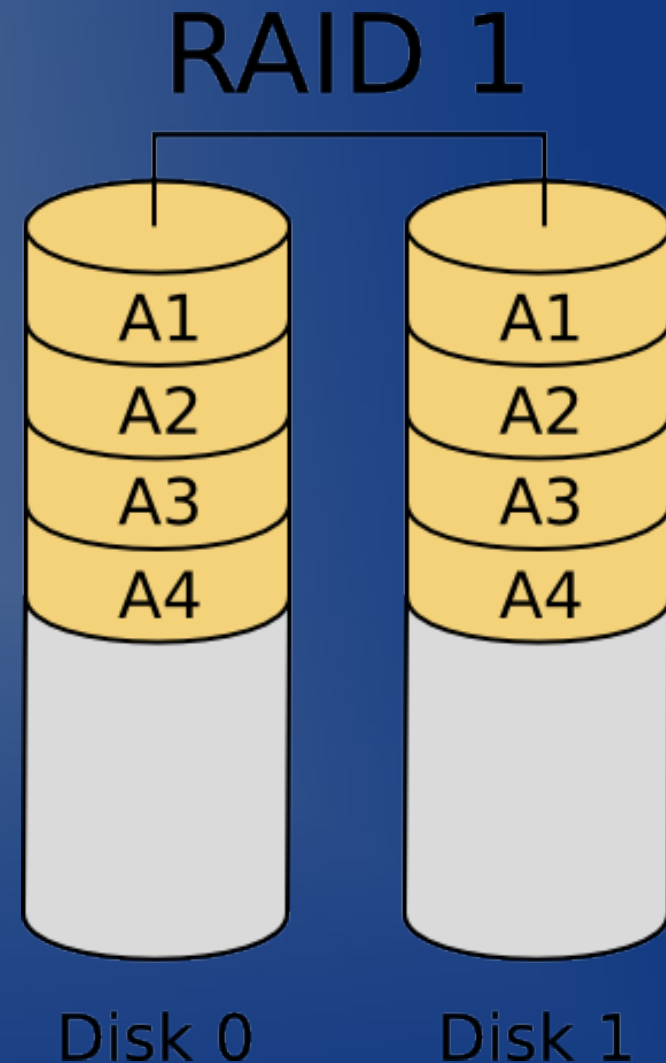
Split Transaction Logs and data.

Use autovacuum (or pay emergency rates)

# RAID 1

Redundancy through  
use of mirror

Increased performance  
(sometimes) through  
shared or partitioned  
reads

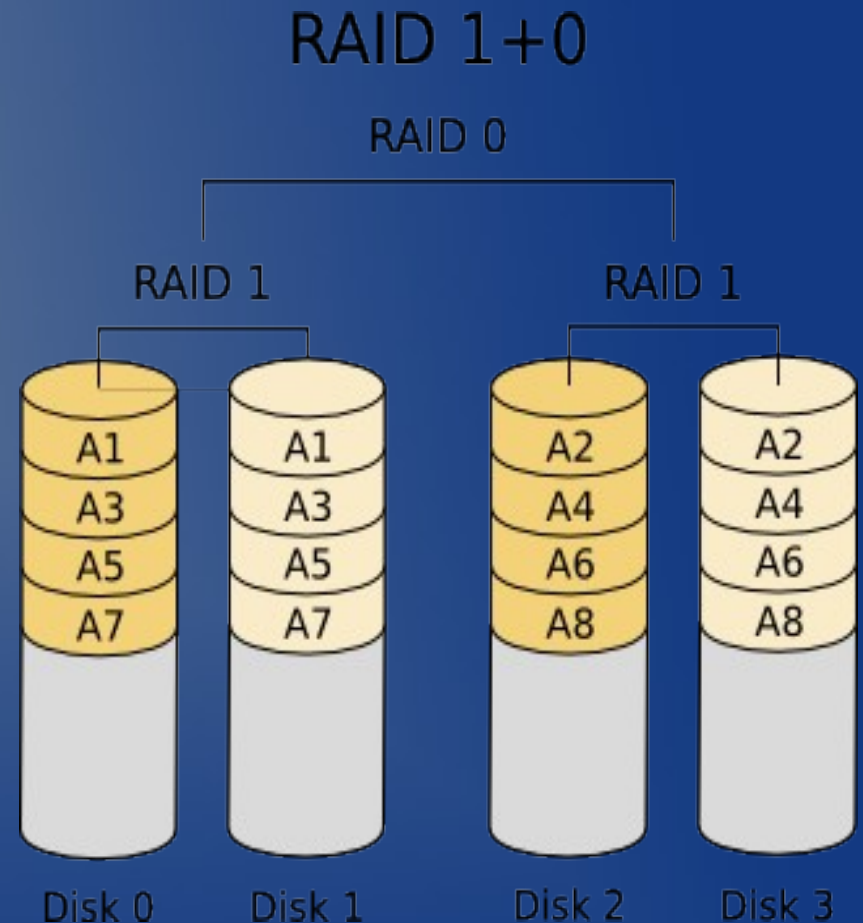


# RAID 1 + 0

Minimum 4 Spindles

Increased performance  
through use of stripe

Increased reliability  
through use of mirror



# Lots of memory

PostgreSQL is efficient and it is possible to run effectively in as little as 256Mb of memory.

Memory is cheap, most data sets are less than 4Gb. If you have at least 4Gb your active data set can remain in file and or `shared_buffer` cache.

# SATA

SATA is great. It is cheap. It is reliable. Just make sure you have a BBU (because of large caches) and you use twice as many spindles.

# Two or More

PostgreSQL is processed based. The more cores you have, the better concurrency you will get.

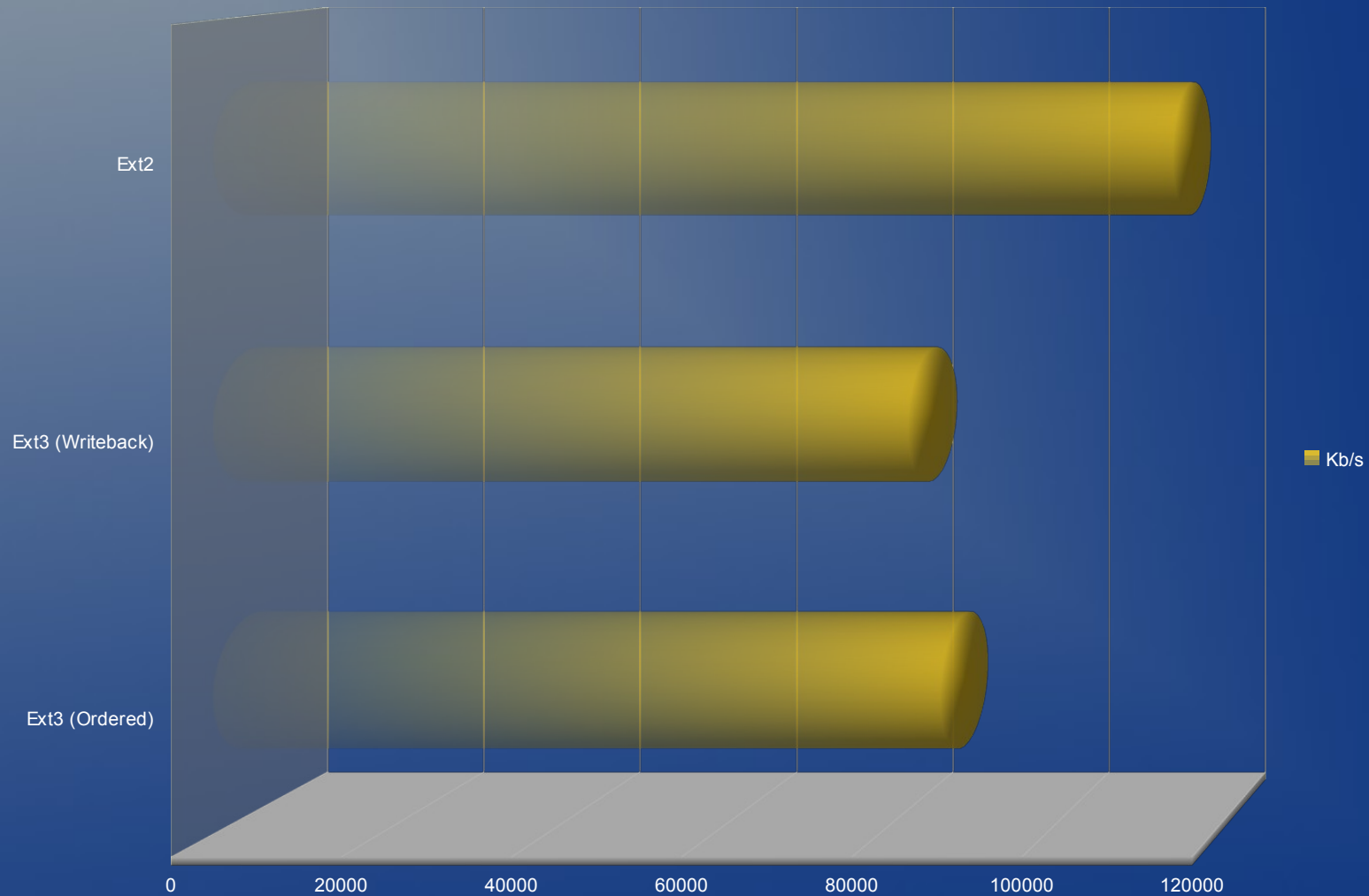
# BBU

## Battery Backup Unit

Used on good RAID cards in case of power outage or sudden crash. Allows for storage of pending writes until the machine comes back on line. A requirement if you are running any kind of CACHE on the RAID or Drives.



# Splitting Transaction Logs



# Autovacuum defaults

```
#autovacuum = on  
#autovacuum_max_workers = 3  
#autovacuum_naptime = 1min  
#autovacuum_vacuum_scale_factor = 0.2  
#autovacuum_analyze_scale_factor = 0.1
```

# Memory

shared\_buffers

work\_mem

maintenance\_work\_mem

effective\_cache\_size

# What are shared\_buffers

The working cache of all hot tuples (and Index entries) within PostgreSQL.

# shared\_buffers

Pre-allocated cache (buffers).

On Linux sysctl.conf – kernel.shmmax

Use 20% of available memory (up to 40%)

Watch out for IO Storms

# What is work\_mem

The working memory available for work operations (sorts) before PostgreSQL will swap.

# work\_mem

Don't set globally (postgresql.conf)

Use per transaction

Can be bad, per query, per connection, per sort

Use EXPLAIN ANALYZE to see if you are overflowing

# Django – SET

```
def my_custom_sql(self):  
    from django.db import connection, transaction  
    cursor = connection.cursor()  
    # Data retrieval operation - no commit required  
    cursor.execute("""  
        SET work_mem TO '%s'""", [self.baz])
```



# Example EXPLAIN ANALYZE

## QUERY PLAN

-----  
Sort (cost=0.02..0.03 rows=1 width=0) (actual time=2270.744..2588.341  
rows=1000000 loops=1)

Sort Key: (generate\_series(1, 1000000))

**Sort Method: external merge Disk: 13696kB**

-> Result (cost=0.00..0.01 rows=1 width=0) (actual  
time=0.006..144.720 rows=1000000 loops=1)

Total runtime: 3009.218 ms  
(5 rows)

# What is maintenance\_work\_mem

The amount of memory (RAM) allowed for maintenance tasks before PostgreSQL swaps. Typical tasks are **ANALYZE, VACUUM, CREATE INDEX, REINDEX**

# maintenance\_work\_mem

Set to a reasonable amount for autovacuum  
Use SET for per session changes such as  
CREATE INDEX

```
SET maintenance_work_mem to '1GB';  
CREATE INDEX foo ON bar(baz);  
RESET maintenance_work_mem;
```

# What is effective\_cache\_size

A pointer for the PostgreSQL planner to hint at how much of the database will be cached. This is not an allocation setting.

# effective\_cache\_size

Take into account shared\_buffers

	total	used	free	shared	buffers	cached
Mem:	6126208	3168356	2957852	0	480884	1258304

% of cached + shared\_buffers = effective\_cache\_size

% depends on workload. Generally between 40% and 70%

# Let's talk IO

log\_checkpoints

checkpoint\_timeout

checkpoint\_completion\_target

checkpoint\_segments

wal\_sync\_method

synchronous\_commit

# log\_checkpoints

By default this is off. Turn on to correlate between checkpoints and spikes in %IOWait from sar.

# checkpoint\_timeout

The amount of time PostgreSQL will wait before it forces a checkpoint. Properly configured it reduces IO utilization. Set to 15 or 20 (minutes). It is affected by:

checkpoint\_segments

checkpoint\_completion\_target



# checkpoint\_completion\_target

This parameter is used to reduce spikes in IO by completing a checkpoint over a period of time.

Do not change this parameter, increase  
checkpoint\_timeout instead.

# checkpoint\_segments

The number of transaction logs that will be utilized before a checkpoint is forced. Each segment is 16 Mb. The default is 3. Use checkpoint\_warning to see if you need more.

Change to at least 10.

Use checkpoint\_warning and logging to get accurate setting.

# wal\_sync\_method

The type of fsync that will be called to flush file modifications to disk. Leave commented to have PostgreSQL figure it out. On Linux it should look like:

```
postgres=# show wal_sync_method ;
wal_sync_method
-----
fdatasync
```

# synchronous\_commit

Specifies whether transaction commit will wait for WAL records to be written to disk before the command returns a "success" indication to the client.

Depends on application. Turn off for faster commits. Low risk of lost commits (but not integrity).

# Let's talk brains

default\_statistics\_target

seq\_page\_cost

random\_page\_cost (to fix)

cpu\_operator\_cost

cpu\_tuple\_cost

# default\_statistics\_target

An arbitrary value used to determine the volume of statistics collected on a relation. The larger the value the longer analyze takes but generally the better the plan. Can be set per column.

# default\_statistics\_target

```
set default_statistics_target to 100;
```

```
pggraph_2_2=# analyze verbose pggraph_indexrollup;
```

```
INFO:  analyzing "aweber_shoggoth.pggraph_indexrollup"
```

```
INFO:  "pggraph_indexrollup": scanned 30000 of 1448084
```

```
pages, containing 1355449 live rows and 0 dead rows; 30000 rows in  
sample, 65426800 estimated total rows
```

```
ANALYZE
```

# default\_statistics\_target

```
set default_statistics_target to 300;
```

```
pggraph_2_2=# analyze verbose pggraph_indexrollup;
```

```
INFO:  analyzing "aweber_shoggoth.pggraph_indexrollup"
```

```
INFO:  "pggraph_indexrollup": scanned 90000 of 1448084
```

```
pages, containing 4066431 live rows and 137 dead rows; 90000 rows in  
sample, 65428152 estimated total rows
```

```
ANALYZE
```

```
pggraph_2_2=#
```



# default\_statistics\_target

How do I know to increase it?

```
Unique (cost=264.65..282.65 rows=100 width=2) (actual time=8.665..12.460
rows=100 loops=1)
  -> Sort (cost=264.65..273.65 rows=3600 width=2) (actual
time=8.664..10.423 rows=3600 loops=1)
    Sort Key: one
    Sort Method: quicksort Memory: 265kB
    -> Seq Scan on bar
(cost=0.00..52.00 rows=52 width=2) (actual time=0.007..1.894 rows=3600
loops=1)
Total runtime: 12.553 ms
```

# Increasing per column

```
ALTER TABLE foo  
  ALTER COLUMN BAR  
    SET STATISTICS 120
```

# seq\_page\_cost

Tells the planner how expensive a sequential scan is. It directly relates to random\_page\_cost.

# random\_page\_cost

Tells the planner the expense of fetching a random page. If using RAID 10, the value should be inverted with seq\_page\_cost (1.0 vs 4.0)

# cpu\_operator\_cost

Sets the planner's estimate of the cost of processing each operator or function executed during a query. The default is 0.0025.

In real world tests, a setting of 0.5 generally provides a better plan. Test using SET in a session.

```
SET cpu_operator_cost TO 0.5;  
EXPLAIN ANALYZE SELECT ...
```

# cpu\_tuple\_cost

Sets the planner's estimate of the cost of processing each row during a query. The default is 0.01.

In real world tests, a setting of 0.5 generally provides a better plan. Test using SET in a session.

```
SET cpu_tuple_cost TO 0.5;  
EXPLAIN ANALYZE SELECT ...
```

# Design

Connection Pooling

Prepared Statements

Functions

Batch Commits

# Connection Pooling

Reduces CPU utilization

Keeps relations hot (in cache)

pgbouncer:

<https://developer.skype.com/SkypeGarage/DbProjects/PgBouncer>



# Prepared Statements

Reduces planning time

Good for recurring and similar transactions

Start: 01:43:08 PM

Insert 1000000 rows 1000 at a time

End: 01:43:18 PM – Time: 10 seconds

Start: 01:45:16 PM

Insert 1000000 prepared rows 1000 at a time

End: 01:45:23 PM – Time: 7 seconds

# Functions

- Reduces processing overhead
  - Multiple round trips
  - Manipulating data inside instead of outside
  - Make sure to test and increase cost as required

# Functions - execution\_cost

A positive number giving the estimated execution cost for the function, in units of `cpu_operator_cost`. If the function returns a set, this is the cost per returned row. If the cost is not specified, 1 unit is assumed for C-language and internal functions, and 100 units for functions in all other languages. Larger values cause the planner to try to avoid evaluating the function more often than necessary.

`result_rows`

# Functions - SET

The SET clause causes the specified configuration parameter to be set to the specified value when the function is entered, and then restored to its prior value when the function exits. SET FROM CURRENT saves the session's current value of the parameter as the value to be applied when the function is entered.

# Functions - ROWS

A positive number giving the estimated number of rows that the planner should expect the function to return. This is only allowed when the function is declared to return a set. The default assumption is 1000 rows.

# CREATE FUNCTION

```
CREATE OR REPLACE FUNCTION RETURN_LOTS(INT) RETURNS SETOF INT AS
$$
    SELECT generate_series(1,$1);
$$
    COST 0.5 ROWS 10000000 SET work_mem TO '5MB'

LANGUAGE 'SQL';
```

# Batch Commits

Reduces commit costs

Increases commit efficiency

Start: 01:45:56 PM

Insert 1000000 rows 1000 at a time

End: 01:46:05 PM – Total: 9 seconds

Start: 01:46:05 PM

Insert 1000000 rows one at a time

End: 01:48:07 PM – Total: 00:2:02

# Questions?

Technical?

Community?